

User Interface - Keypad Scan, PSoC® Style

AN2034

Author: Dave Van Ess

Associated Project: Yes

Associated Part Family: CY8C20x34, CY8C21x23, CY8C21x34

CY8C23x33, CY8C24x23A, CY8C24x94

CY8C27x43, CY8C29x66

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Software Version: PSoC Designer™ 5.0 Associated Application Notes: AN2354

Application Note Abstract

X–Y matrix keypads are an inexpensive interface enabling interaction with microcontroller-based products. This application note shows how the PSoC® microcontroller's unique I/O structure can build a keypad scan routine that is fast, uses minimal RAM resources, and operates in a polled or interrupt mode. A function callable by either 'C' or assembly language is also presented.

Introduction

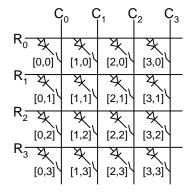
An X-Y keypad enables use of N column lines and M row lines to detect switch closures for N * M switches. For this application note, a keypad is defined as an X-Y matrix where only one key is pressed at a time, as opposed to a keyboard where simultaneous key closures are the norm ([Ctrl] [Shift] [Delete]). This keypad definition is valid for telephones, calculators, security entry kiosks, or other products where only one key is pressed at a time.

This application relies on PSoC General Purpose Input Output (GPIO).

Rows and Columns

This application note uses C columns and R rows. Figure 1 shows an example of such a keypad:

Figure 1. 4-Column by 4-Row Keypad



Closure of switch [i, j] (column i, row j) enables current flow from row j to column i. This keypad requires only eight connections to the MCU. The 16 diodes can detect multiple key closures. However, because PSoC microcontroller reduces the cost of external components such as op-amps, filters, and DACs, using 16 diodes is not a good idea.

Well known techniques have been developed to detect multiple key presses without diodes. Figure 2 shows the keypad without diodes.

Figure 2. The Keypad you can Afford

	C	o C) ₁ C	c_2) ₃
R ₀	[0,0]	[1,0]	[2,0]	[3,0]	
R ₁	[0,1]	[1,1]	[2,1]	[3,1]	
R ₂ -	[0,2]	[1,2]	[2,2]	[3,2]	
R ₃	[0,3]	[1,3]	[2,3]	[3,3]	

The standard algorithm for reading a keypad is to individually drive each row and sample the status of all columns. Correctly combining all this information enables detection of at least two simultaneous switch closures. The hardware cost is less but there is software overhead required to scan all four rows, read column status, and condense this information into an answer.

To develop a keypad scan that is low in both hardware and software resources, limit its operation to single-key presses.

PSOC General Purpose IO

To detect single-switch closures, use the following algorithm:

- Drive all rows simultaneously and read the columns.
- Drive all columns simultaneously and read the rows.
- Condense this data to determine switch-closure status.

The structure of the general purpose of each pin simplifies this bipolar use of rows and columns. Each pin has a digital driver that can be set up to be:

- Strong drive to V_{DD} and a pull down resistor to V_{ss}.
- Strong drive to V_{ss} and a pull up resistor to V_{DD}.
- High Impedance
- Strong drive to either V_{DD} or V_{ss}

Use the pull down drive mode in the keypad scan routine, as shown in Figure 3. Because it is also the default condition for each pin at initial startup, there is no requirement for special port configuration.

Figure 3. GPIO Set For Pull Down Mode

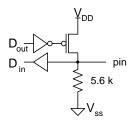
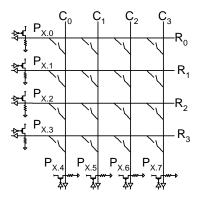


Figure 4 shows implementation of a keypad scan.

Figure 4. PSoC Architecture for Keypad Scan



For this, press example switch [2,1] (column 2, row 1). The algorithm reads the keypad in six steps.

- Output b00001111 to the port. This drives all the rows high, leaving the columns passively pulled down.
- Read the port. The driven pins 0 through 3 remain high and because the switch [2,1] is closed, pin 6 is now high. The value read is b01001111.
- Output b11110000 to the port. This drives all the columns high, leaving the rows passively pulled down.
- Read the port. The driven pins 4 through 7 remain high and because the switch [2,1] is closed, pin 1 is now high. The value read is b11110010.
- "Anding" the result of step 2 and step 4 results in the answer b01000010.
- The upper 4 bits decode as column 2 and the lower 4 bits decode as row 1. This is a match with the closed switch.

A subroutine that implements this algorithm is shown in example Code 1. It enables the reading of a four row by four column keypad connected to port 1. It is found in "Keypad.asm," located in the project file associated with this application note. This subroutine uses eight instructions, 15 bytes of program memory bytes, and 57-CPU cycles.

Code 1: Subroutine for Reading Keypad

```
Keypad.asm
  This routine reads a 4 column by 4 row
  keypad on port1. The status of key
  closures is returned in A.
         P1.4 P1.5 P1.6 P1.7
         C0 C1 C2 C3
 P1.0 R0 --+---+--
           P1.1 R1 --+-
             P1.2 R2 --+---
          ;
 P1.3 R3 --+---+-
export bReadKeypad export bReadKeypad include "m8c.inc"
 bReadKeypad:
_bReadKeypad:
  mov reg[PRT1DR], f0h ;drive columns
  mov
       X,SP
  mov A, reg[PRT1DR]
                        ;read rows
  mov reg[PRT1DR], 0fh ;drive rows
                  ;store row info on stack
  push A
                        ;Read Columns
  mov A, reg[PRT1DR]
  and [X], A
                         ; combine them
  pop A
ret
```

The C header shown in example Code 2 can be found in "Keypad.h." It makes the subroutine shown in example Code 1 a 'C' callable function.

Code 2. C Header Example

```
// Create a pragma to support
// proper argument and return
// value passing
#pragma fastcall bReadKeypad
extern BYTE bReadKeypad(void);
```

Using the Output

The output of the function bReadKeypad is a single byte that shows the status of key closure of the keypad. It is translated and decoded as follows:

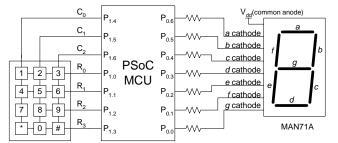
- No bits are set if no key is pressed.
- A single bit in the upper nibble and a single bit in the lower nibble are set for a single-key press.
- Any other condition is a multiple-key closure and is defined as not valid.

These rules can be decoded with discrete conditional code that breaks up the byte into two nibbles to determine row and column information. Use this information to determine which key, if any, was pressed. This results in a complex set of rules and is tedious.

Another scheme is to use a lookup table to decode this data. The advantage is that the table stores the formatted data. Different programmers could be working on the same project and each use their own table to decode the keypad when they are required to read it.

The project file associated with this application note uses such a table to decode the key closures. A block diagram of the project is shown in Figure 5.

Figure 5. Block Diagram for the Keypad Project



For this project, set the drives for the port 1 pins to the Pull Down mode (default) and the port 0 pins to either the Pull Down or Strong mode.

The keypad is scanned and the appropriate bits are set on the output port to turn on the desired LED segments. This display is a single digit 7-segment common anode LED display. Any particular segment is lit when its cathode is pulled low. As an example, all output pins low result in an "8" being displayed. All output pins high result in a blank display.

Example Code 3 is the main function that implements the design in Figure 5 on page 3.

Code 3. Keypad Project Implemented

```
;-----
; This program reads the keypad
; at port 1 and control the LEDs
; on port0.
; Copyright (c)
; Cypress Semiconductor 2002-2008.
; All Rights Reserved.
include "m8c.inc"
export main
main:
loop:
  call bReadKeypad
  index KeyToLED
        reg[PRTODR], A
  mov
        loop
  qmr
  ret
xxh: equ 30h ;illegal character "E"
.Literal
KeyToLED:
  db 7fh,xxh,xxh,xxh,xxh,xxh,xxh
  db xxh,xxh,xxh,xxh,xxh,xxh,xxh
  db xxh,4fh,4ch,xxh,0fh,xxh,xxh,xxh
  db 78h,xxh,xxh,xxh,xxh,xxh,xxh
  db xxh,12h,24h,xxh,00h,xxh,xxh,xxh
  db 01h,xxh,xxh,xxh,xxh,xxh,xxh
  db xxh,xxh,xxh,xxh,xxh,xxh,xxh
  db xxh,xxh,xxh,xxh,xxh,xxh,xxh
  db xxh,06h,20h,xxh,04h,xxh,xxh,xxh
  db 18h,xxh,xxh,xxh,xxh,xxh,xxh
  db xxh,xxh,xxh,xxh,xxh,xxh,xxh
  db xxh,xxh,xxh,xxh,xxh,xxh,xxh
  db xxh,xxh,xxh,xxh,xxh,xxh,xxh
  db xxh,xxh,xxh,xxh,xxh,xxh,xxh
  db xxh,xxh,xxh,xxh,xxh,xxh,xxh
  db xxh,xxh,xxh,xxh,xxh,xxh,xxh
.EndLiteral
```

The ".Literal" and ".EndLiteral" macros are used to disable the compiler's code compression algorithms so that the literal information in the KeyToLED table is not compressed.

This code can also be found in the project associated with this application note. The table makes for a simple translation from key press to displayed character. Develop your own table for a particular application.

Debouncing

Debouncing solutions are specific to each system's program structure and switch selection. However, debouncing is simple; it requires two sequential keypad scans to agree for the scan to be valid.

Interrupt Driven Keypad Scans

Execute the following steps for an interrupt-instigated key scan:

- Configure the lower four (row) pins to be I/O interrupts on a rising edge.
- Set the upper four (column) pins high.
- Write an interrupt handler that calls bReadKeypad.
- Enable the GPIO interrupt.

The row pins stay low until a key press connects a column to a row causing an interrupt.

One possible application is a security gate keypad reader where the MCU is placed in the power-saving-sleep mode. When someone presses a key, the processor wakes up and decodes the user's input. After finishing, the MCU can put itself back to sleep. For more information about sleep mode operation please refer to application note AN2354.

Summary

The PSoC microcontroller GPIO structure is ideal for fast keypad scans. The "index" instruction enables simple translation key presses to a user's desired format. The GPIO to interrupts makes for an interrupt-driven keypad scan.

Document History

Document Title: User Interface – Keypad Scan, PSoC® Style

Document Number: 001-40409

Revision	ECN	Orig. of Change	Submission Date	Description of Change	
**	1532004	OGNE	10/02/07	New publication of existing application note.	
*A	2640952	JVY	01/20/09	Updated content. Added part numbers CY8C20x34, CY8C21x23, CY8C21x34, CY8C23x33, CY8C24x23A, CY8C24x94, CY8C27x43, and CY8C29x66.	

In March of 2007, Cypress recataloged all of its Application Notes using a new documentation number and revision code. This new documentation number and revision code (001-xxxxx, beginning with rev. **), located in the footer of the document, will be used in all subsequent revisions.

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